

Safely Measuring Tor

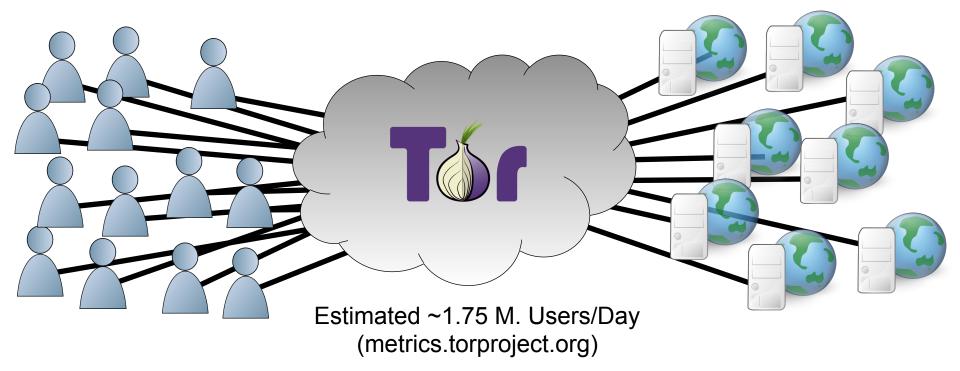
"Safely Measuring Tor", Rob Jansen and Aaron Johnson, In the *Proceedings of the 23rd ACM Conference on* Computer and Communication Security (CCS 2016).

Rob Jansen

U.S. Naval Research Laboratory Center for High Assurance Computer Systems 23rd Conference on Computer and Communication Security Hofburg Imperial Palace, Vienna, Austria October 27th, 2016



Talk Overview

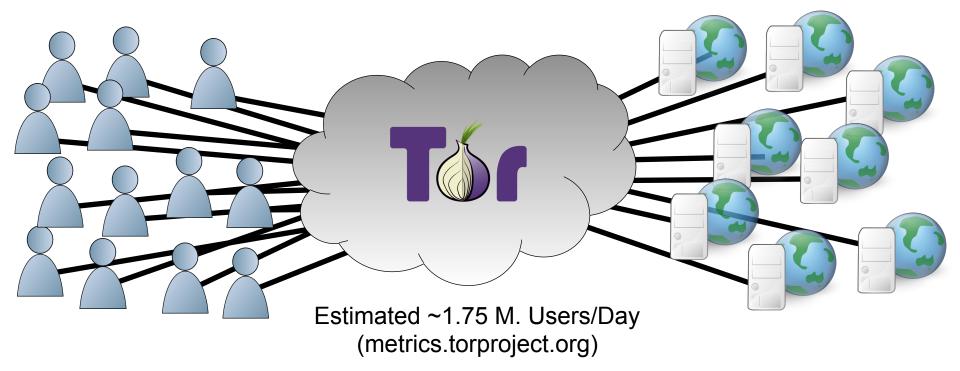


Tor: an anonymous communication, censorship resistant, privacy-enhancing communication system

How is Tor being used? being misused? performing?



Talk Overview



Tor: an anonymous communication, censorship resistant, privacy-enhancing communication system

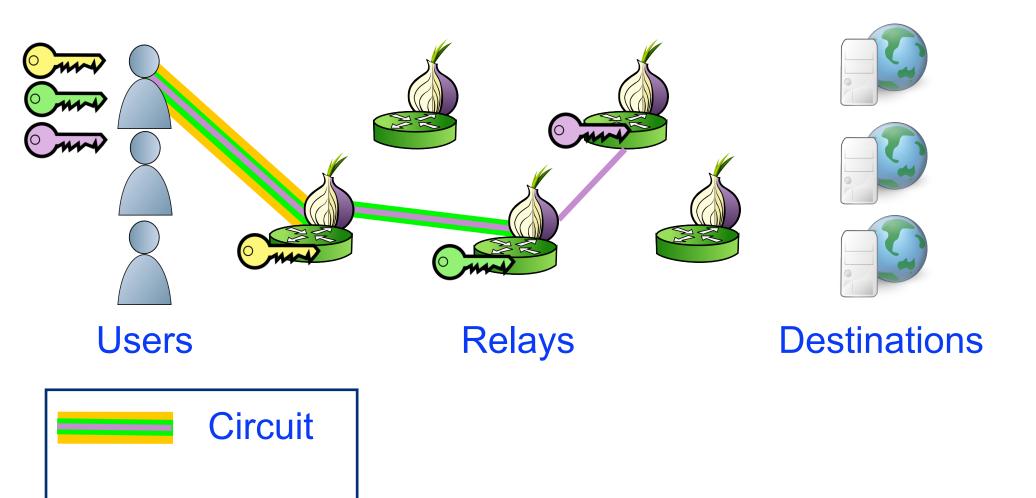
- How is Tor being used? being misused? performing?
- Objective: To safely gather Tor network usage statistics
- Approach: Use distributed measurement, secure multiparty computation, and differential privacy

Background and Motivation

- How Tor works
- Why measurements are needed and what to measure
- Measurement challenges

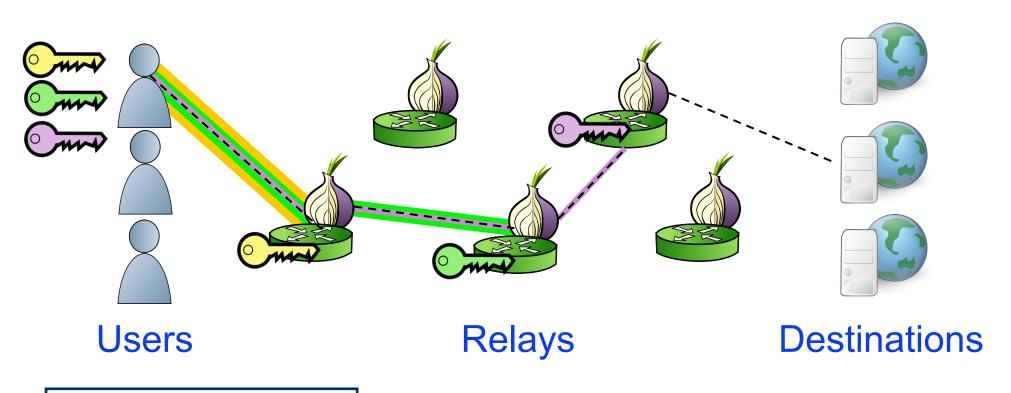


Background: Onion Routing

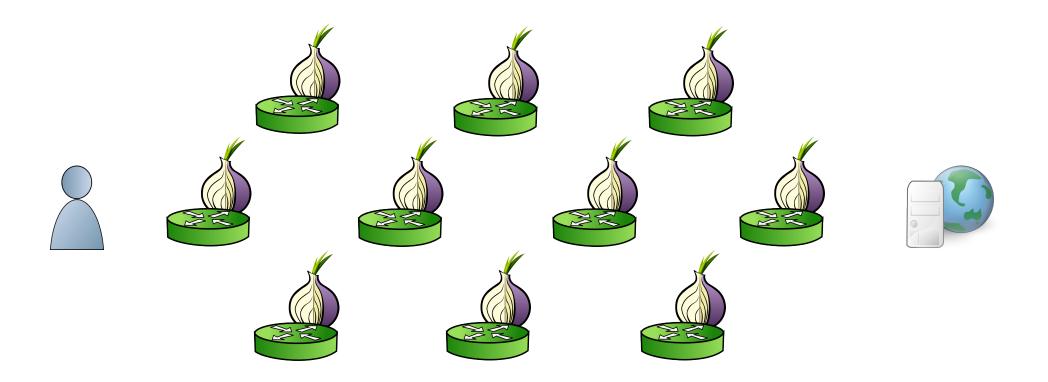




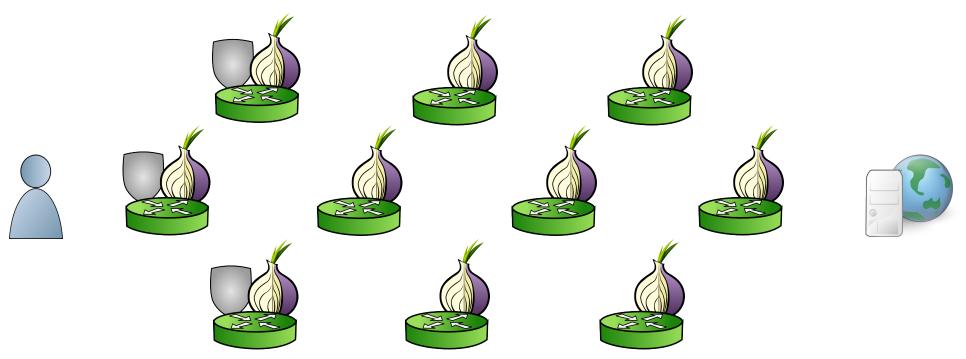
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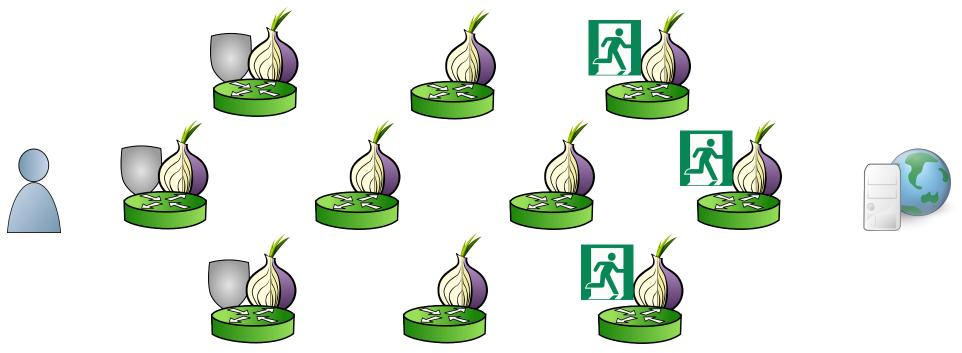






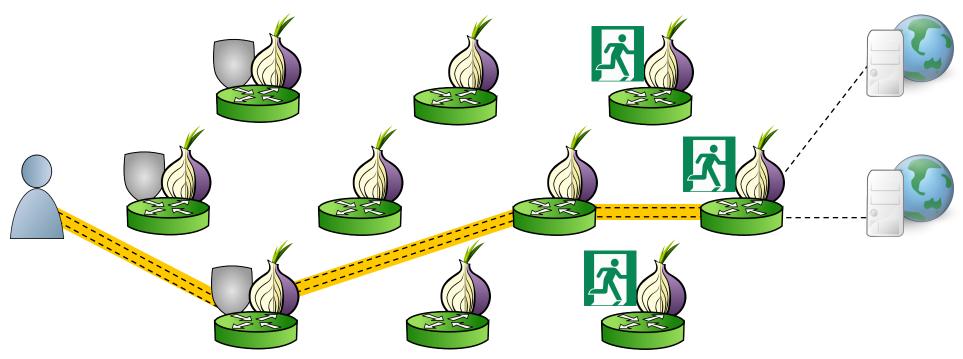
1. Clients begin all circuits with a selected guard





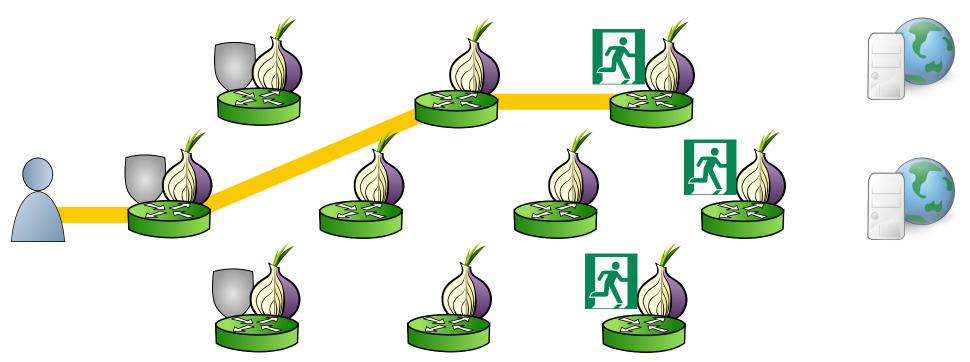
- 1. Clients begin all circuits with a selected guard
- Relays define individual exit policies





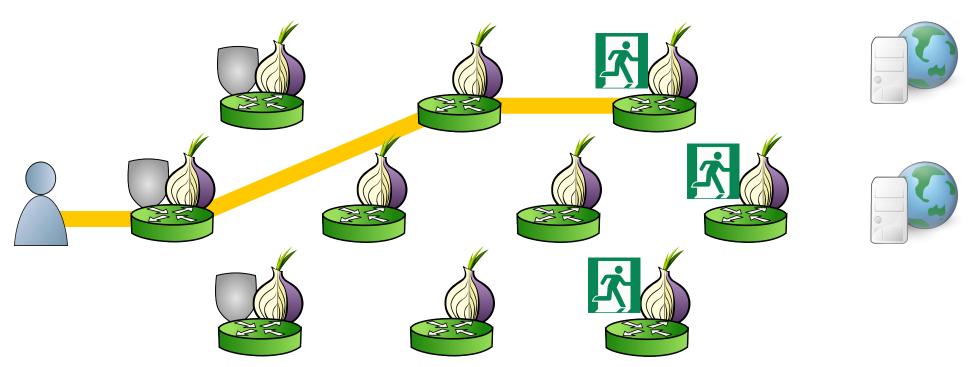
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- Relays define individual exit policies
- 3. Clients multiplex streams over a circuit





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- 1. Clients begin all circuits with a selected guard
- 2. Relays define individual exit policies
- 3. Clients multiplex streams over a circuit
- 4. New circuits replace existing ones periodically
- 5. Clients randomly choose relays, weighted by bandwidth



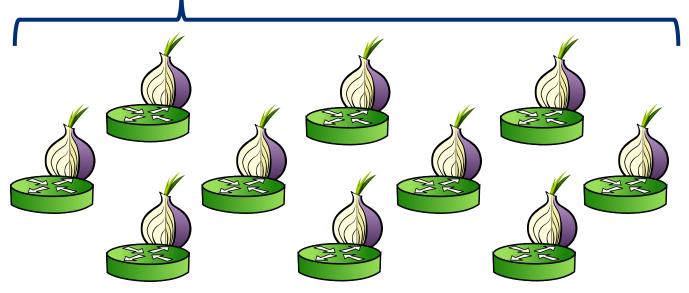
Background: Directory Authorities

Directory Authorities



Hourly network consensus by majority vote

- Relay info (IPs, pub keys, bandwidths, etc.)
- Parameters (performance thresholds, etc.)





Motivation: Why Measure Tor?

Why are Tor network measurements needed?

- To understand usage behaviors to focus effort and resources
- To understand network protocols and calibrate parameters
- To inform policy discussion



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"Tor metrics are the ammunition that lets Tor and other security advocates argue for a more private and secure Internet from a position of data, rather than just dogma or perspective."

> Bruce Schneier (June 1, 2016) (metrics.torproject.org)



Motivation: Previous Measurement Studies

Previous work collected, stored, and manually analyzed sensitive data

- McCoy et. al. (PETS 2008): tcpdump of first 150 bytes of packet (including 96 payload)
- Chaabane et. al. (NSS 2010): customized DPI software



face both civil and criminal penalties for a research project in which they snooped on users of the Tor anonymous proxy network. Should federal prosecutors take inte

July 24, 2008 by Chris Soghoian 9:40 AM PDT



Motivation: Measurement Challenges



Some Existing Measurements

Data Published	Privacy Techniques	Unsafe	Inaccurate
Relay BW available	Test measurements		*
Relay BW used	Aggregated ~ 4 hours	×	
Total # daily users	Inferred (consensus fetches)		*
# users per country	Aggregated ~ 24 hours, rounded, opt-in	*	
Exit traffic per port	Aggregated ~ 24 hours, opt-in	*	



Motivation: Measurement Challenges



Some Existing Measurements

Safety concerns:

- Per-relay outputs
- Data stored locally
- No privacy proofs

Data Published	Privacy Techniques	Unsafe	Inaccurate
Relay BW available	Test measurements		*
Relay BW used	Aggregated ~ 4 hours	*	
Total # daily users	Inferred (consensus fetches)		*
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Motivation: Measurement Challenges



Some Existing Measurements

Accuracy concerns:

- Per-relay noise
- Opt-in, limited vantage points

Data Published	Privacy Techniques	Unsafe	Inaccurate
Relay BW available	Test measurements		*
Relay BW used	Aggregated ~ 4 hours	×	
Total # daily users	Inferred (consensus fetches)		*
# users per country	Aggregated ~ 24 hours, rounded, opt-in	*	
Exit traffic per port	Aggregated ~ 24 hours, opt-in	*	



Motivation: Missing Measurements

Many useful statistics are not collected for safety

Users

Total number of unique users at any time, how long they stay online, how often they join and leave, usage behavior

Relays

Total bandwidth capacity, congestion and queuing delays, circuit and other failures, denial of service and other attacks

Destinations

Popular destinations, popular applications, effects of DNS, properties of traffic (bytes and connections per page, etc.)

The PrivCount Measurement System

- PrivCount system architecture
- Distributed measurement and aggregation protocol
- Secure computation and private output



PrivCount: Overview

Privacy-preserving counting system

- Consumes various new event types from Tor
 - Circuit end events
 - Stream end events
 - Connection end events
- Counts various statistics from event information, e.g.:
 - Total number of circuits, streams, connections
 - Data volume per circuit, stream
 - Number of unique users
- Based on PrivEx-S2 protocol of Elahi et. al. (CCS 2014)



PrivCount: Overview

Security goals for safer Tor measurements



- Forward privacy
 - The adversary cannot learn the state of the measurement before time of compromise
- Differential privacy
 - Prevents confirmation of the actions of a specific user given the output
- Secure aggregation
 - Securely aggregates safe statistics across all measurement nodes
 - Only the safe, aggregated measurement results are released



PrivCount: Architecture

Data Collectors (DCs)

- Collect events
- Increment counters



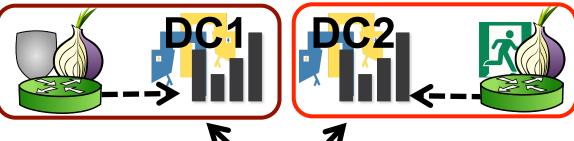




PrivCount: Architecture

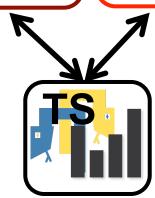
Data Collectors (DCs)

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Tally Server (TS)

- Central, untrusted proxy
- Collection facilitator





PrivCount: Architecture

Data Collectors (DCs)

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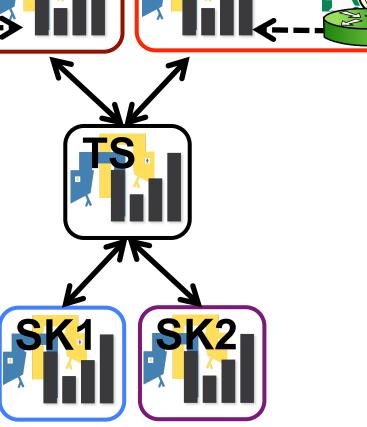


Tally Server (TS)

- Central, untrusted proxy
- Collection facilitator

Share Keepers (SKs)

 Stores DC secrets, sum for aggregation



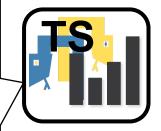


PrivCount: Initialization



Create deployment document

- Privacy parameters ε and δ
- Sensitivity for each statistic (max change due to single client)
- Noise weight ω
 (relative noise added by each DC)



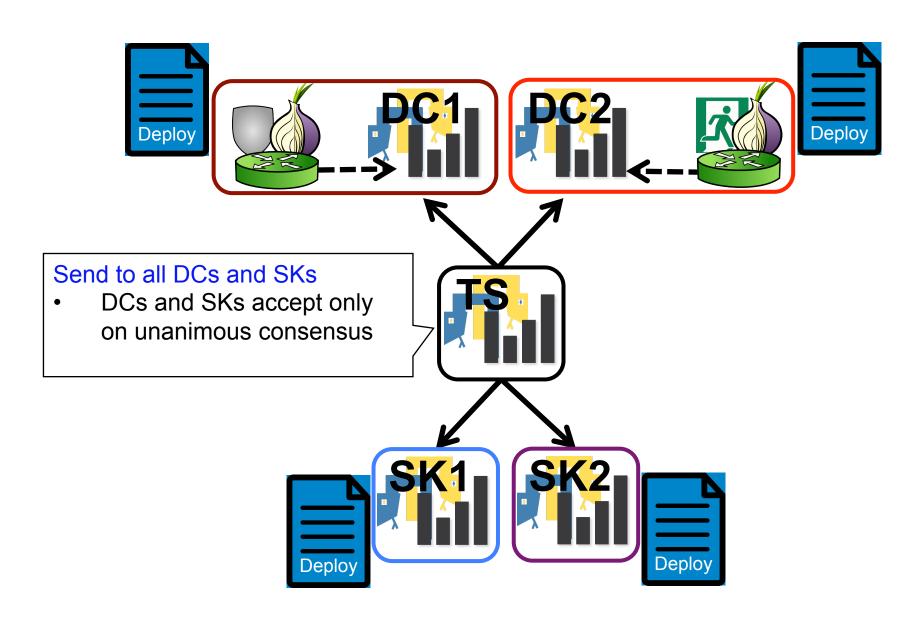








PrivCount: Initialization





PrivCount: Configuration



Create configuration document

- Collection start and end times
- Statistics to collect
- Estimated value for each statistic (maximize relative per-statistic accuracy while providing (ε, δ)differential privacy)



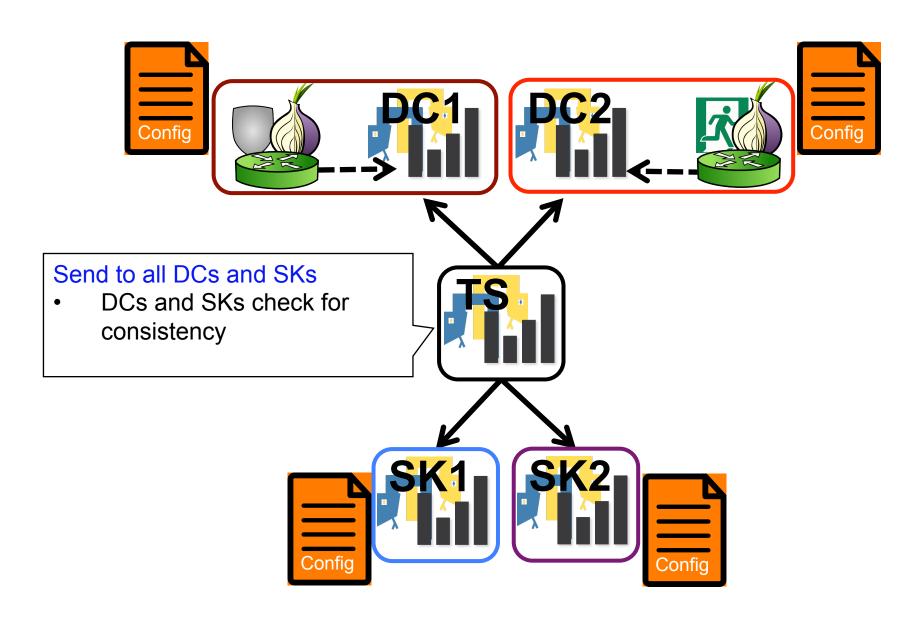








PrivCount: Configuration





PrivCount: Execution - Setup



Generate noise for each counter

- $N \sim Normal(0,\omega\sigma) \mod q$
- Contributes to differential privacy of the outputs









PrivCount: Execution - Setup



Generate random share for each SK

- $S \sim Uniform(\{0, ..., q-1\})$
- "Blinds" the actual counts for forward privacy at the DCs

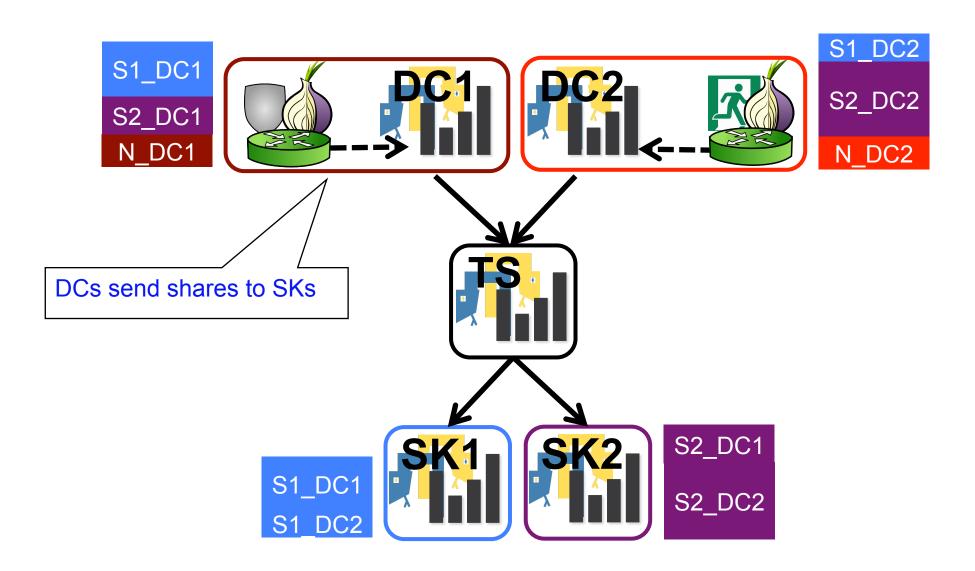








PrivCount: Execution - Setup





PrivCount: Collection

C_DC1

S1_DC1

S2_DC1 N_DC1





C_DC2

S1_DC2

S2_DC2

DCs collect events and increment counters



S1_DC1



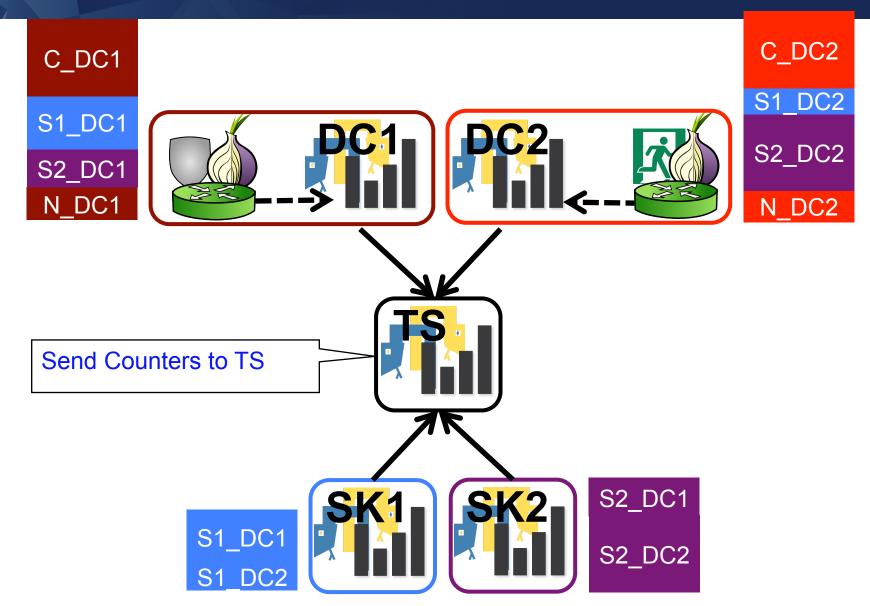


S2_DC1

S2_DC2

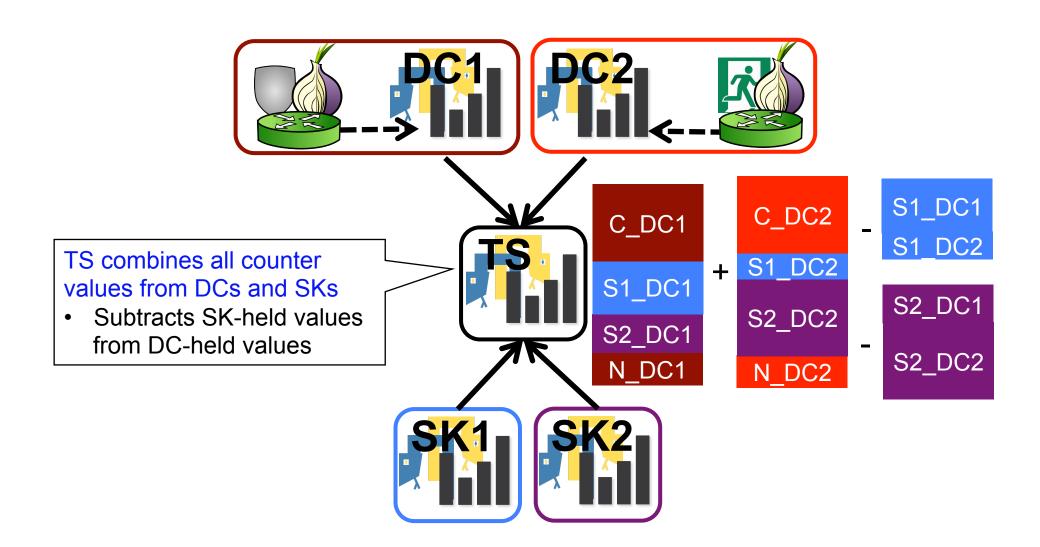


PrivCount: Collection



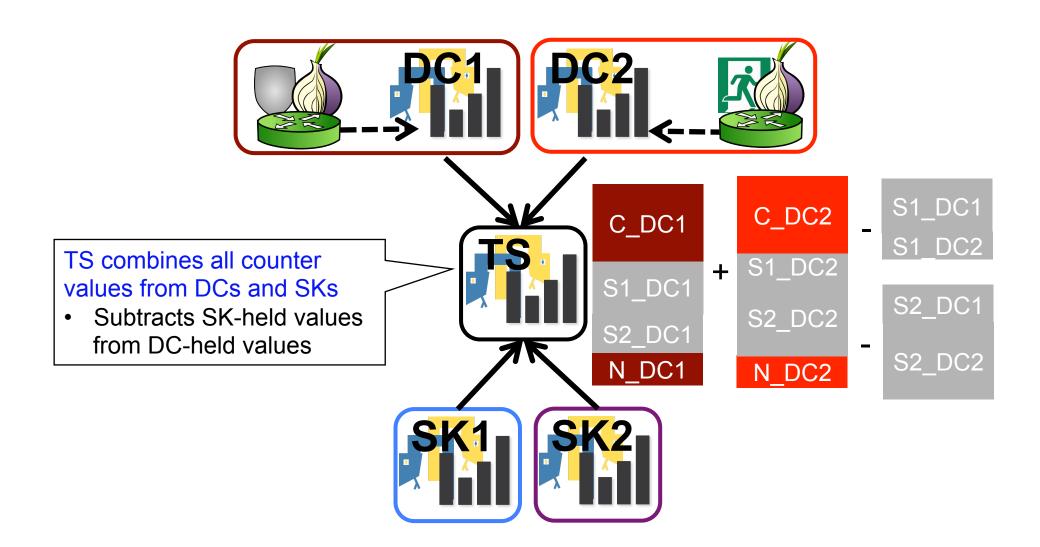


PrivCount: Aggregation



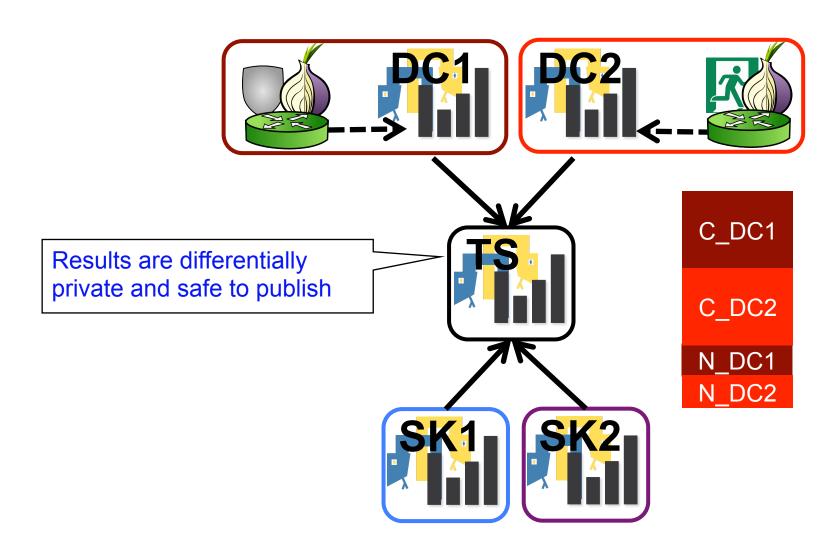


PrivCount: Aggregation





PrivCount: Aggregation





Recall: Security Properties

- Forward privacy
 - The adversary cannot learn the state of the measurement before time of compromise
- Differential privacy
 - Prevents confirmation of the actions of a specific user given the output
- Secure aggregation
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C_DC1

S1_DC1

S2_DC1

N_DC1



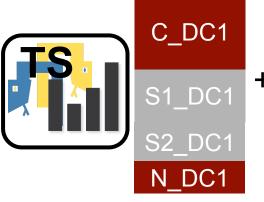


C_DC2

S1 DC2

S2_DC2

N_DC2



C_DC2

S1_DC2

S2_DC2

N_DC2

S1_DC1 S1_DC2

S2_DC1

S2_DC2

S1_DC1 S1_DC2





S2_DC1

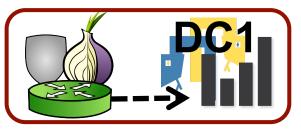
S2_DC2



C_DC1

S1_DC1

S2_DC1 N_DC1





C_DC2

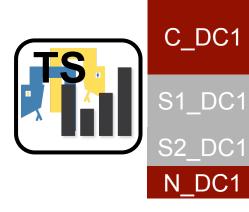
S1 DC2

S2_DC2

N DC2

Forward Privacy

Nothing learned from counter before time of compromise as long as 1 SK is honest



C_DC2 S1 DC2 S2 DC2 N DC2

S1 DC1 S1 DC2

S2 DC1

S2_DC2

S1_DC1





S2_DC1 S2_DC2



C_DC1

S1_DC1

S2_DC1 N_DC1





C_DC2

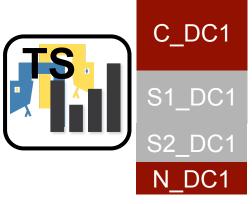
S1 DC2

S2_DC2

N DC2

Differential Privacy

Enough noise is added as long as a tunable subset of DCs are honest



C_DC2 S1 DC2 S2 DC2

N DC2

S1 DC1 S1 DC2

S2 DC1

S2_DC2

S1_DC1





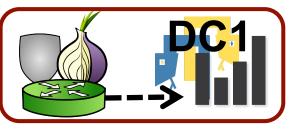
S2_DC1 S2_DC2



C_DC1

S1_DC1

S2_DC1 N_DC1





C_DC2

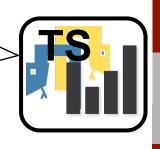
S1 DC2

S2_DC2

N_DC2

Secure Aggregation

Count+noise is added securely
the TS only learns the aggregated sum



C_DC1

S1_DC1 S2_DC1

N_DC1

C_DC2

S1_DC2

S2_DC2

N_DC2

S1_DC1

S2 DC1

S1 DC2

S2_DC2

S1_DC1 S1_DC2





S2_DC1

S2_DC2



C_DC1

S1_DC1

S2_DC1 N_DC1





C_DC2

S1 DC2

S2_DC2

N DC2

See paper for more details and for security and privacy proofs



C_DC1 S1 DC1 S2 DC1 N_DC1

C_DC2 S1 DC2

S2 DC2

N DC2

S1 DC1 S1 DC2

S2 DC1

S2_DC2

S1_DC1





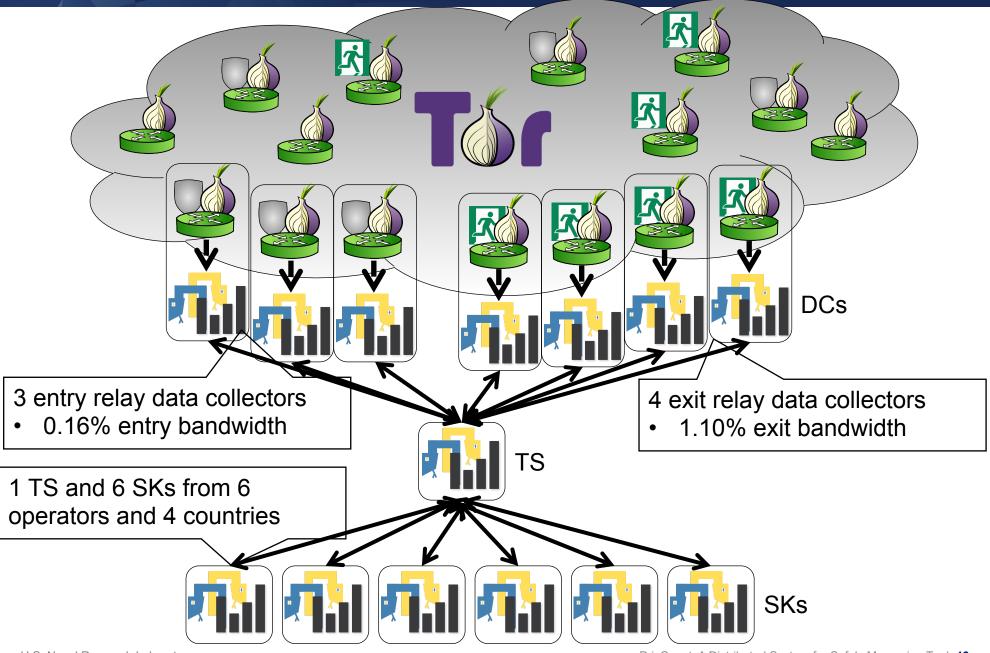
S2_DC1 S2_DC2

Deployment and Measurement Results

- Configuring and running Tor relays
- "Exploratory" measurements using various exit policies
- "In-depth" measurements of most popular usage
- Network-wide measurement inference



Deploying PrivCount





Collection Phases

Exploratory phases

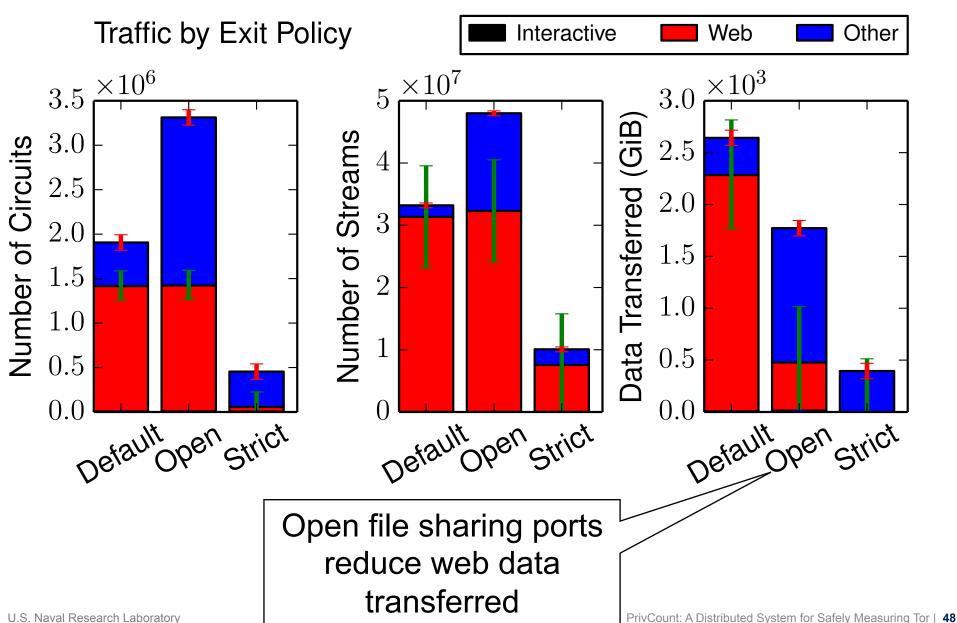
- Explore various exit policies (strict, default, open)
- Explore various applications (web, interactive, other)
- Gather only totals (circuits, streams, bytes)
- Use Tor metrics to estimate input parameters
- Run for 1 day, iterate

In-depth phases

- Focus on most popular exit policy and applications
- Gather totals and histograms
- Use exploratory results to estimate input parameters
- Run for 4 days for client stats, 21 days for exit stats

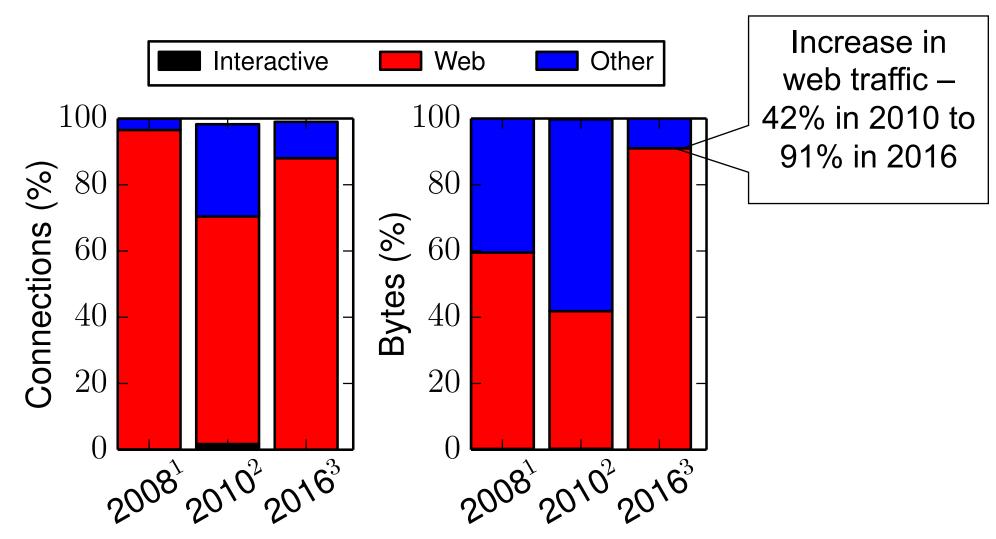


Results: Exit Policies





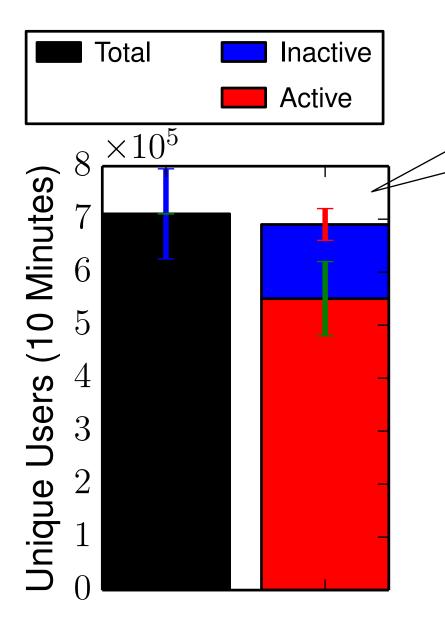
Results: Amount and Types of Traffic



[1] PETS 2008, McCoy... [2] NSS 2010, Chaabane... [3] CCS 2016, Jansen...



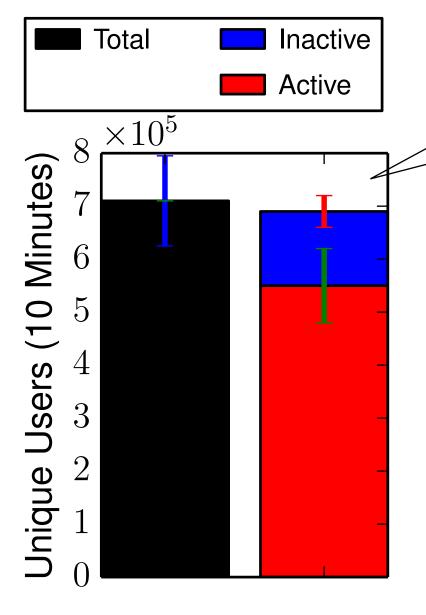
Results: Number of Unique Users



710k total users 550k (77%) active users In an average 10 mins.



Results: Number of Unique Users

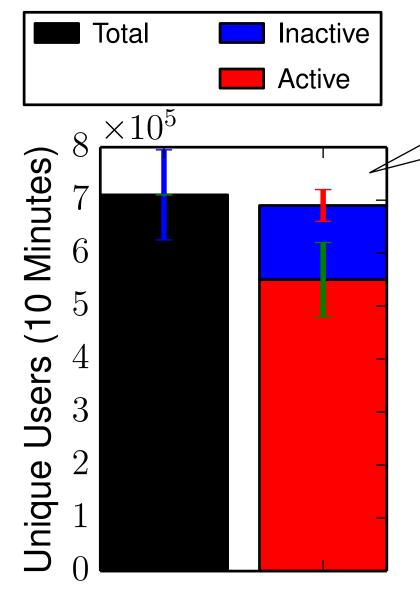


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~800k – ~1.6m average concurrent users (Tor Browser update pings – https://tor-metrics.shinyapps.io/ webstats2/)



Results: Number of Unique Users



710k total users 550k (77%) active users In an average 10 mins.

~800k – ~1.6m average concurrent users (Tor Browser update pings – https://tor-metrics.shinyapps.io/ webstats2/)

~1.75m daily users (Consensus downloads https://metrics.torproject.org)



Results: Traffic Modeling Statistics

More results in the paper!

Table 11: Distributions of Tor network activity from histogram-counter in-depth exit statistics

Statistic		Bin Ranges and Count Distribution (with ± 95% CI)							
Active Circuit Life Time (s)		[1, 480):	57%±44%	[480, 720):	45%±42%	[720, 1200):	0%±33%	[1200, ∞):	0%±35%
Streams Per Circuit	Total Web Other	[1, 3): [1, 3): [1, 3):	46%±43% 36%±37% 78%±15%	[3, 7): [3, 7): [3, 7):	38%±41% 22%±33% 10%±9%	(7, 15): (7, 15): (7, 15):	31%±40% 13%±31% 0%±8%	[15, ∞): [15, ∞): [15, ∞):	9%±37% 3%±28% 2%±8%
Client-bound Bytes Per Stream	Total Web Other	[1, 2048): [1, 2048): [1, 2048):	60%±40% 33%±33% 56%±21%	(2048, 16384): (2048, 16384): (2048, 16384):	38%±35% 37%±34% 9%±15%	(16384, 65536): (16384, 65536): (16384, 65536):	32%±33% 5%±26% 8%±15%	[65536, ∞): [65536, ∞): [65536, ∞):	6%±26% 0%±24% 11%±15%
Server-bound Bytes Per Stream	Total Web Other	[1, 512): [1, 512): [1, 512):	57%±39% 41%±35% 40%±19%	[512, 1024): [512, 1024): [512, 1024):	25%±31% 36%±34% 6%±14%	[1024, 4096): [1024, 4096): [1024, 4096):	38%±34% 23%±30% 15%±16%	[4096, ∞): [4096, ∞): [4096, ∞):	0%±24% 2%±25% 1%±14%
Bytes Per Stream Ratio	Total Web Other	(-∞, -1): (-∞, -1): (-∞, -1):	80%±45% 70%±42% 45%±20%	[-1, 1): [-1, 1): [-1, 1):	25%±31% 15%±28% 14%±16%	(1, ∞): (1, ∞): (1, ∞):	0%±21% 0%±21% 12%±15%		
Inter-stream Creation Time (s)	Total Web Other	[0, 1): [0, 1): [0, 1):	87%±47% 68%±41% 16%±16%	(1, 5): (1, 5): (1, 5):	16%±29% 8%±27% 10%±15%	(5, 10): (5, 10): (5, 10):	1%±25% 13%±28% 3%±14%	[10, ∞): [10, ∞): [10, ∞):	0%±23% 14%±28% 12%±15%



Conclusion

PrivCount

- Distributed measurement system using secret sharing
- Safer Tor measurement study
- Open source: https://github.com/privcount

Future measurement plans

- Network traffic to create realistic traffic models
- Onion services to improve reliability and scalability
- Better techniques for cardinality (e.g., # unique users)
- Detecting denial of service attacks and other misbehavior

Contact

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Questions



PrivCount vs PrivEx

How does PrivCount enhance PrivEx

- Multi-phase iterative measurement
- Expanded privacy notion that simultaneously handles multiple types of measurements
- Optimal allocation of the ε privacy budget across multiple statistics
- Composable security definition and proof
- More capable and reliable tool
- Supports over 30 different types of Tor statistics
- Resilience to node failures and reboots
- Simpler configuration and setup



Privacy

Parameters for (ε, δ) -differential privacy

- $\varepsilon = 0.3$: same as used by Tor onion service stats
- $\delta = 10^{-3}$: upper bound on prob. of choosing noise value that violates ε-differential privacy
- DCs on 3 machines, add 3x noise

User action bounds

Action	Bound	
Simultaneous open entry connections	1	
Entry connection open time	24 hours	
New entry connections	12	
New circuits	146	
New streams	30,000	
Data sent or received	10 MiB	